

FINITE MATHEMATICS

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for Business, Economics,
Life Sciences, and
Social Sciences



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Chapter 6

Linear Programming: The Simplex Method

Section 3

The Dual Problem:
Minimization with
Problem Constraints of
the Form $\geq$

Learning Objectives for Section 6.3



Dual Problem: Minimization with Problem Constraints of the Form $\geq$

- The student will be able to formulate the dual problem.
- The student will be able to solve minimization problems.
- The student will be able to solve applications of the dual such as the transportation problem.
- The student will be able to summarize problem types and solution methods.

Dual Problem: Minimization With Problem Constraints of the Form $\geq$

Associated with each minimization problem with $\geq$ constraints is a maximization problem called the **dual problem**. The dual problem will be illustrated through an example. We wish to minimize the objective function C subject to certain constraints:

$$C = 16x_1 + 9x_2 + 21x_3$$

$$x_1 + x_2 + 3x_3 \geq 12$$

$$2x_1 + x_2 + x_3 \geq 16$$

$$x_1, x_2, x_3 \geq 0$$

Initial Matrix



We start with an initial matrix A which corresponds to the problem constraints:

$$A = \begin{bmatrix} 1 & 1 & 3 & 12 \\ 2 & 1 & 1 & 16 \\ 16 & 9 & 21 & 1 \end{bmatrix}$$

Transpose of Matrix A



To find the transpose of matrix A , interchange the rows and columns so that the first row of A is now the first column of A transpose. The transpose is used in formulating the dual problem to follow.

$$A^T = \begin{bmatrix} 1 & 2 & 16 \\ 1 & 1 & 9 \\ 3 & 1 & 21 \\ 12 & 16 & 1 \end{bmatrix}$$

Dual of the Minimization Problem

The dual of the minimization problem is the following maximization problem:

Maximize P under the constraints:

$$P = 12y_1 + 16y_2$$

$$y_1 + 2y_2 \leq 16$$

$$y_1 + y_2 \leq 9$$

$$3y_1 + y_2 \leq 21$$

$$y_1, y_2 \geq 0$$

Formation of the Dual Problem



Given a minimization problem with $\geq$ problem constraints,

1. Use the coefficients and constants in the problem constraints and the objective function to form a matrix A with the coefficients of the objective function in the last row.
2. Interchange the rows and columns of matrix A to form the matrix A^T , the transpose of A .
3. Use the rows of A^T to form a maximization problem with $\leq$ problem constraints.

Theorem 1: Fundamental Principle of Duality



A minimization problem has a solution if and only if its dual problem has a solution. If a solution exists, then the optimal value of the minimization problem is the same as the optimum value of the dual problem.

Forming the Dual Problem



We transform the inequalities into equalities by adding the slack variables x_1, x_2, x_3 :

$$y_1 + 2y_2 + x_1 = 16$$

$$y_1 + y_2 + x_2 = 9$$

$$3y_1 + y_2 + x_3 = 21$$

$$-12y_1 - 16y_2 + p = 0$$

Form the Simplex Tableau for the Dual Problem

The first pivot element is 2 (in red) because it is located in the column with the smallest negative number at the bottom (-16), and when divided into the rightmost constants yields the smallest quotient ($16/2=8$)

$$\begin{array}{c|cccccc} & y_1 & y_2 & x_1 & x_2 & x_3 & P \\ \hline x_1 & 1 & \textcolor{red}{2} & 1 & 0 & 0 & 16 \\ x_2 & 1 & 1 & 0 & 1 & 0 & 9 \\ x_3 & 3 & 1 & 0 & 0 & 1 & 21 \\ P & -12 & -16 & 0 & 0 & 0 & 0 \end{array}$$

Simplex Process

Divide row 1 by the pivot element (2) and change the entering variable to y_2

$$\begin{array}{c} y_2 \\ x_2 \\ x_3 \\ P \end{array} \begin{pmatrix} y_1 & y_2 & x_1 & x_2 & x_3 & P \\ .5 & 1 & .5 & 0 & 0 & 8 \\ 1 & 1 & 0 & 1 & 0 & 9 \\ 3 & 1 & 0 & 0 & 1 & 21 \\ -12 & -16 & 0 & 0 & 0 & 0 \end{pmatrix}$$

Simplex Process (continued)

Perform row operations to get zeros in the column containing the pivot element. Result is shown below.

Identify the next pivot element (0.5) (in red)

	y_1	y_2	x_1	x_2	x_3	P
y_2	.5	1	.5	0	0	8
New pivot element → x_2	.5	0	-.5	1	0	1
x_3	2.5	0	-.5	0	1	13
Pivot element located in this column because of negative indicator → P	-4	0	8	0	0	128

Simplex Process (continued)

- Variable y_1 becomes new entering variable
- Divide row 2 by 0.5 to obtain a 1 in the pivot position (circled.)

$$\begin{array}{c}
 y_2 \\
 y_1 \\
 x_3 \\
 P
 \end{array}
 \left(\begin{array}{cccccc}
 y_1 & y_2 & x_1 & x_2 & x_3 & P \\
 .5 & 1 & .5 & 0 & 0 & 8 \\
 1 & 0 & -1 & 2 & 0 & 2 \\
 2.5 & 0 & -.5 & 0 & 1 & 13 \\
 -4 & 0 & 8 & 0 & 0 & 128
 \end{array} \right)$$

Simplex Process (continued)

Get zeros in the column containing the new pivot element.

$$\begin{array}{c}
 y_1 \quad y_2 \quad x_1 \quad x_2 \quad x_3 \quad P \\
 y_2 \left(\begin{array}{ccccc|c}
 0 & 1 & 1 & -1 & 0 & 8 \\
 1 & 0 & -1 & 2 & 0 & 2 \\
 0 & 0 & 2 & -5 & 1 & 8 \\
 0 & 0 & 4 & 8 & 0 & 136
 \end{array} \right)
 \end{array}$$

We have now obtained the optimal solution since none of the bottom row indicators are negative.

Solution of the Linear Programming Problem

- **Solution:** An optimal solution to a minimization problem can always be obtained from the bottom row of the final simplex tableau for the dual problem.
- Minimum of P is 136, which is also the maximum of the dual problem. It occurs at $x_1 = 4$, $x_2 = 8$, $x_3 = 0$

	y_1	y_2	x_1	x_2	x_3	P
y_2	0	1	1	-1	0	8
y_1	1	0	-1	2	0	2
x_3	0	0	2	-5	1	8
P	0	0	4	8	0	136

Solution of a Minimization Problem



Given a minimization problem with non-negative coefficients in the objective function,

1. Write all problem constraints as $\geq$ inequalities. (This may introduce negative numbers on the right side of some problem constraints.)
2. Form the dual problem.
3. Write the initial system of the dual problem, using the variables from the minimization problem as slack variables.

Solution of a Minimization Problem



4. Use the simplex method to solve the dual problem.
5. Read the solution of the minimization problem from the bottom row of the final simplex tableau in step 4.

Note: If the dual problem has no optimal solution, the minimization problem has no optimal solution.

Application: Transportation Problem

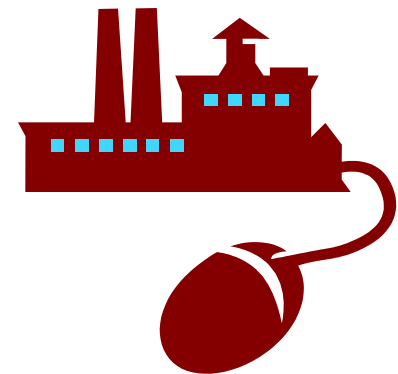


One of the first applications of linear programming was to the problem of minimizing the cost of transporting materials. Problems of this type are referred to as **transportation problems**.

Example: A computer manufacturing company has two assembly plants, plant A and plant B, and two distribution outlets, outlet I and outlet II. Plant A can assemble at most 700 computers a month, and plant B can assemble at most 900 computers a month. Outlet I must have at least 500 computers a month, and outlet II must have at least 1,000 computers a month.

Transportation Problem (continued)

Transportation costs for shipping one computer from each plant to each outlet are as follows: \$6 from plant A to outlet I; \$5 from plant A to outlet II; \$4 from plant B to outlet I; \$8 from plant B to outlet II. Find a shipping schedule that will minimize the total cost of shipping the computers from the assembly plants to the distribution outlets. What is the minimum cost?



Transportation Problem (continued)



Solution: To form a shipping schedule, we must decide how many computers to ship from either plant to either outlet. This will involve 4 decision variables:

x_1 = number of computers shipped from plant A to outlet I

x_2 = number of computers shipped from plant A to outlet II

x_3 = number of computers shipped from plant B to outlet I

x_4 = number of computers shipped from plant B to outlet II

Transportation Problem (continued)



Constraints are as follows:

$$x_1 + x_2 \leq 700 \text{ Available from A}$$

$$x_3 + x_4 \leq 900 \text{ Available from B}$$

$$x_1 + x_3 \geq 500 \text{ Required at I}$$

$$x_2 + x_4 \geq 1,000 \text{ Required at II}$$

Total shipping charges are:

$$C = 6x_1 + 5x_2 + 4x_3 + 8x_4$$

Transportation Problem (continued)

Thus, we must solve the following linear programming problem:

$$\text{Minimize } C = 6x_1 + 5x_2 + 4x_3 + 8x_4$$

subject to

$$x_1 + x_2 \leq 700 \text{ Available from A}$$

$$x_3 + x_4 \leq 900 \text{ Available from B}$$

$$x_1 + x_3 \geq 500 \text{ Required at I}$$

$$x_2 + x_4 \geq 1,000 \text{ Required at II}$$

Before we can solve this problem, we must multiply the first two constraints by -1 so that all are of the $\geq$ type.

Transportation Problem (continued)

The problem can now be stated as:

$$\text{Minimize } C = 6x_1 + 5x_2 + 4x_3 + 8x_4$$

subject to

$$-x_1 - x_2 \geq -700$$

$$-x_3 - x_4 \geq -900$$

$$x_1 + x_3 \geq 500$$

$$x_2 + x_4 \geq 1,000$$

$$x_1, x_2, x_3, x_4 \geq 0$$

Transportation Problem (continued)

$$A = \left[\begin{array}{cccc|c} -1 & -1 & 0 & 0 & -700 \\ 0 & 0 & -1 & -1 & -900 \\ 1 & 0 & 1 & 0 & 500 \\ 0 & 1 & 0 & 1 & 1,000 \\ \hline 6 & 5 & 4 & 8 & 1 \end{array} \right]$$

$$A^T = \left[\begin{array}{cccc|c} -1 & 0 & 1 & 0 & 6 \\ -1 & 0 & 0 & 1 & 5 \\ 0 & -1 & 1 & 0 & 4 \\ 0 & -1 & 0 & 1 & 8 \\ \hline -700 & -900 & 500 & 1,000 & 1 \end{array} \right]$$

Transportation Problem (continued)

The dual problem is;

$$\text{Maximize } P = -700y_1 - 900y_2 + 500y_3 + 1,000y_4$$

subject to

$$-y_1 + y_3 \leq 6$$

$$-y_1 + y_4 \leq 5$$

$$-y_2 + y_3 \leq 4$$

$$-y_2 + y_4 \leq 8$$

$$y_1, y_2, y_3, y_4 \geq 0$$

Transportation Problem (continued)

Introduce slack variables x_1 , x_2 , x_3 , and x_4 to form the initial system for the dual:

$$-y_1 \quad + y_3 \quad + x_1 \quad = 6$$

$$-y_1 \quad + y_4 \quad + x_2 \quad = 5$$

$$-y_2 + y_3 \quad + x_3 \quad = 4$$

$$-y_2 \quad + y_4 \quad + x_4 \quad = 8$$

$$-700y_1 - 900y_2 + 500y_3 + 1,000y_4 \quad + P \quad = 0$$

Transportation Problem Solution

If we form the simplex tableau for this initial system and solve, we find that the shipping schedule that minimizes the shipping charges is 0 from plant A to outlet I, 700 from plant A to outlet II, 500 from plant B to outlet I, and 300 from plant B to outlet II. The total shipping cost is \$7,900.



Summary of Problem Types and Simplex Solution Methods

Problem	Constraints of Type	Right-Side Constants	Coefficients of Objective Function	Method of Solution
Maximization	$\leq$	Nonnegative	Any real number	Simplex Method
Minimization	$\geq$	Any real number	Nonnegative	Form dual and solve by simplex method